

Set-Based Modeling in LocalSolver 6.0

www.localsolver.com

Who we are



Bouygues, one of the French largest corporation, €33 bn in revenues http://www.bouygues.com

Innovation24

Operations Research subsidiary of Bouygues 20 years of practice and research

http://www.innovation24.fr

LocalSolver

Mathematical optimization solver commercialized by Innovation 24 http://www.localsolver.com

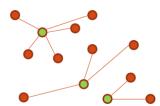


LocalSolver

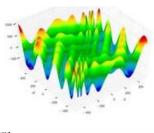
Hybrid math optimization solver

LocalSolver = LS + CP/SAT + LP/MIP + NLP

For combinatorial, numerical, or mixed-variable optimization



Suited for large-scale non-convex optimization



Quality solutions in seconds without tuning





free trial with support – free for academics – renting licenses from 590 €/month – perpetual licenses from 9,900 €

Motivations

Modeling approaches for the *Traveling Salesman Problem*



Mixed-Integer Programming

With an **exponential number** of constraints



Minimise $\sum c_{ij} x_{ij}$

Variant with O(n²) variables and constraints

Conventional Formulation (C) (Dantzig, Fulkerson and Johnson (1954))

 $\forall i \in N$

 $\sum x_{ij} = 1$

 $\forall j \in N$ $\sum_{i} x_{ij} = 1$

 $\sum x_{ij} \leq |M| - 1$ $\forall M \subset N \text{ such that } \{1\} \notin M, |M| \geq 2$

SINGLE COMMODITY FLOW (F1) (Gavish and Graves (1978))

Both constraints are retained but we also introduce (continuous) variables:

 y_{ij} = 'Flow' in an arc (i,j) $i \neq j$

and constraints:

 $y_{ij} \leq (n-1)x_{ij}$

$$\forall i,j \in N, i \neq j$$

→ Iterative procedure to add subtour elimination constraints

$$\sum_{\substack{j\\j\neq 1}} y_{1j} = n - 1$$

$$\sum_{i} y_{ij} - \sum_{k} y_{jk} = 1 \qquad \forall j \in \mathbb{N} - \{1\}$$

$$i \neq i \qquad i \neq k$$



LocalSolver 5.0

A compact model based on positions

```
/* Declares the optimization model. */
□function model() {
    // x[i][j] equal to 1 if city j is ith visited city in the tour
    x[0..nbCities-1][0..nbCities-1] <- bool();
    // one city per position i
    for [i in 0..nbCities-1] constraint sum[j in 0..nbCities-1](x[i][j]) == 1;
    // one position per city i
    for [j in 0..nbCities-1] constraint sum[j in 0..nbCities-1](x[j][j]) == 1;
    //city[i] is the city at position i in the tour
    city[i in 0..nbCities-1] \leftarrow sum[j in 0..nbCities-1](j*x[i][j]);
    // Distance of the arc reaching the ith city
    distance[0] <- distanceWeight[city[nbCities - 1]][city[0]];</pre>
    distance[i in 1..nbCities - 1] <- distanceWeight[city[i - 1]][city[i]];</pre>
    // Minimize the total distance
    obj <- sum[i in 0..nbCities - 1](distance[i]);</pre>
     minimize obj;
```

- → A polynomial-size model (not an iterative procedure)
- → No artificial variables and constraints
- → But still based on binary decisions



Natural Modeling

As a permutation

The Traveling Salesman Problem (TSP)

TSP: Given a list of cities and their pairwise distances, find a shortest possible tour that visits each city exactly once.

Objective: find a permutation $a_1,...,a_n$ of the cities that minimizes

$$d(a_1, a_2) + d(a_2, a_3) + \dots + d(a_{n-1}, a_n) + d(a_n, a_1)$$

where d(i, j) is the distance between cities i and j



An optimal TSP tour through Germany's 15 largest cities

Reference modeling

Garey & Johnson

[ND22] TRAVELING SALESMAN

INSTANCE: Set C of m cities, distance $d(c_i, c_j) \in Z^+$ for each pair of cities $c_i, c_j \in C$, positive integer B.

QUESTION: Is there a tour of C having length B or less, i.e., a permutation $\langle c_{\pi(1)}, c_{\pi(2)}, \ldots, c_{\pi(m)} \rangle$ of C such that

$$\left(\sum_{i=1}^{m-1} d(c_{\pi(i)}, c_{\pi(i+1)})\right) + d(c_{\pi(m)}, c_{\pi(1)}) \leq B ?$$



Set-based modeling

Innovative modeling concepts for routing & scheduling problems



List Variables

Structured decisional operator list(n)

- Order a subset of values in domain {0, ..., n-1}
- Each value is unique in the list

Classical operators to interact with "list"

- count(u): number of values selected in the list
- at(u,i) or u[i]: value at index i in the list
- indexOf(u,v): index of value v in the list
- contains(u,v): equivalent to "indexOf(u,v) != -1"
- disjoint(u1, u2, ..., uk): true if u1, u2, ..., uk are pairwise disjoint
- partition(u1, u2, ..., uk): true if u1, u2, ..., uk induce a partition of {0, ..., n-1}



Traveling salesman

TSP: Given a list of cities and their pairwise distances, find a shortest possible tour that visits each city exactly once.

Objective: find a permutation $a_1,...,a_n$ of the cities that minimizes

$$d(a_1, a_2) + d(a_2, a_3) + \dots + d(a_{n-1}, a_n) + d(a_n, a_1)$$





where d(i, j) is the distance between cities i and j

Why not a single line model?

constraint TSP (graph);



Real-world models do not fit in tight frameworks

Time[A, B, T] =
$$\frac{2(\alpha_{x}^{2} + \alpha_{y}^{2})}{-2(\alpha_{x}\beta_{x} + \alpha_{y}\beta_{y}) + \sqrt{4(\alpha_{x}\beta_{x} + \alpha_{y}\beta_{y})^{2} - 4(\alpha_{x}^{2} + \alpha_{y}^{2})(\beta_{x}^{2} + \beta_{y}^{2} - V^{2})}} + T$$

With α_x , β_x , α_y , β_y function of A,B and T

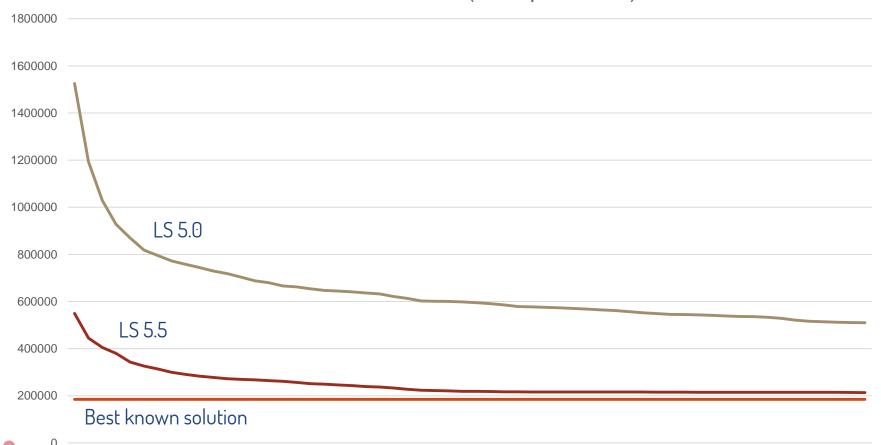
constraint TSP(graph);



Kinetic TSP

Performance?

TSP: real-life 200-client instance LocalSolver 5.0 vs 5.5 (with operator *list*)



Vehicle routing

	TSP	VRP
Normal	Count(x)=N	partition(x[1K])
Prize-collecting	maximize sum()	disjoint(x[1K])



CVRP benchmarks

CVRP - instances A

- 32 to 80 clients, 10 trucks max.
- 27 instances
- 5 minutes of running time
- LS binary: almost infeasible
- LS list: 1 % avg. opt. gap

CVRP - instances X100-500

- 100 to 500 clients, 138 trucks max.
- 67 instances
- 5 minutes of running time
- LS binary: almost infeasible
- LS list: 9 % avg. opt. gap



CVRPTW benchmarks

CVRPTW - instances Solomon R100

- 101 to 112 clients, 19 trucks max.
- 13 instances
- 5 minutes of running time
- LS binary: N/A
- LS list: 3 % avg. opt. gap

CVRPTW - instances Solomon R200

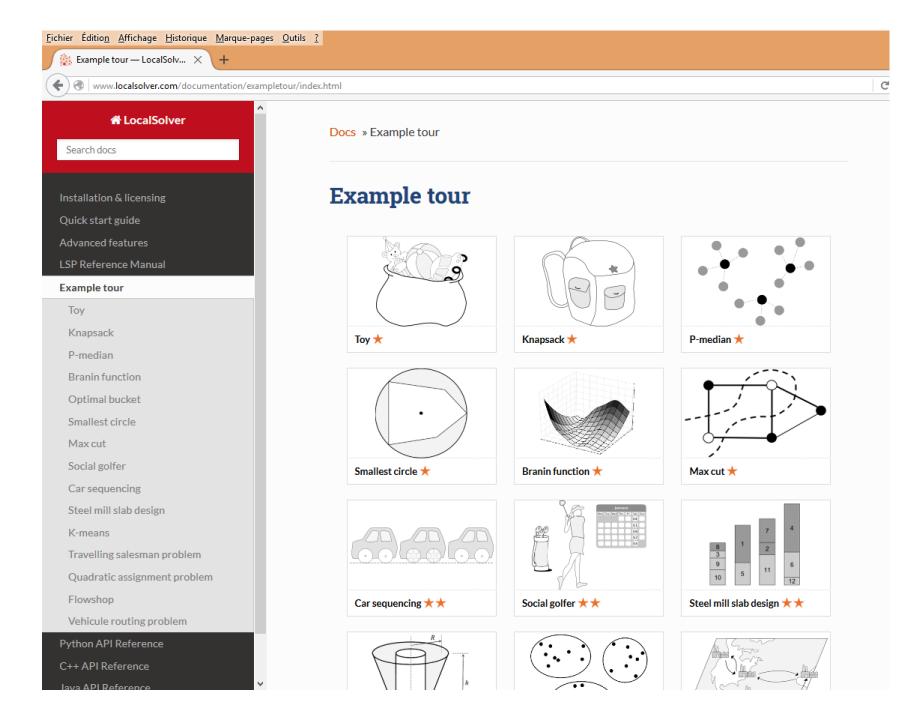
- 201 to 208 clients, 4 trucks max.
- 8 instances
- 5 minutes of running time
- LS binary: N/A
- LS list: 8 % avg. opt. gap



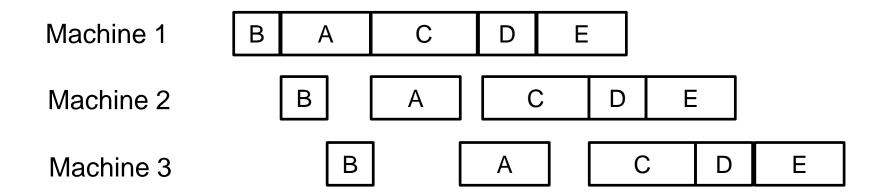
Beyond routing problems

Scheduling, planning, sequencing





Flow-shop scheduling



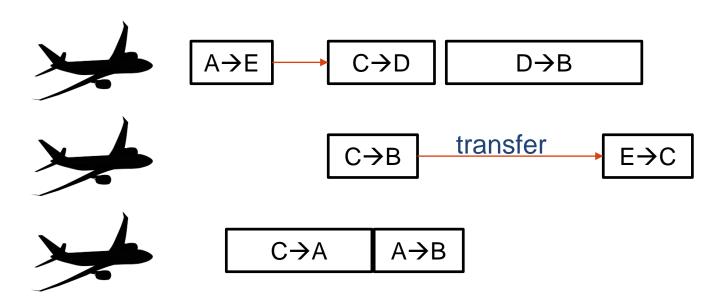
Since we are looking for a permutation of jobs the model is straightforward with a single list variable



Planning

Flights to plane assignments





A solution is a partition of flights into K lists (one per plane)

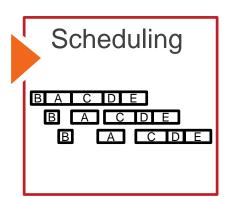
The goal is to minimize the total transfer times

Conclusion

List Variables are a first step towards set-based modeling in LocalSolver

This higher level of modeling yields simple and compact models producing high quality solutions for









Any other sequencing problem

